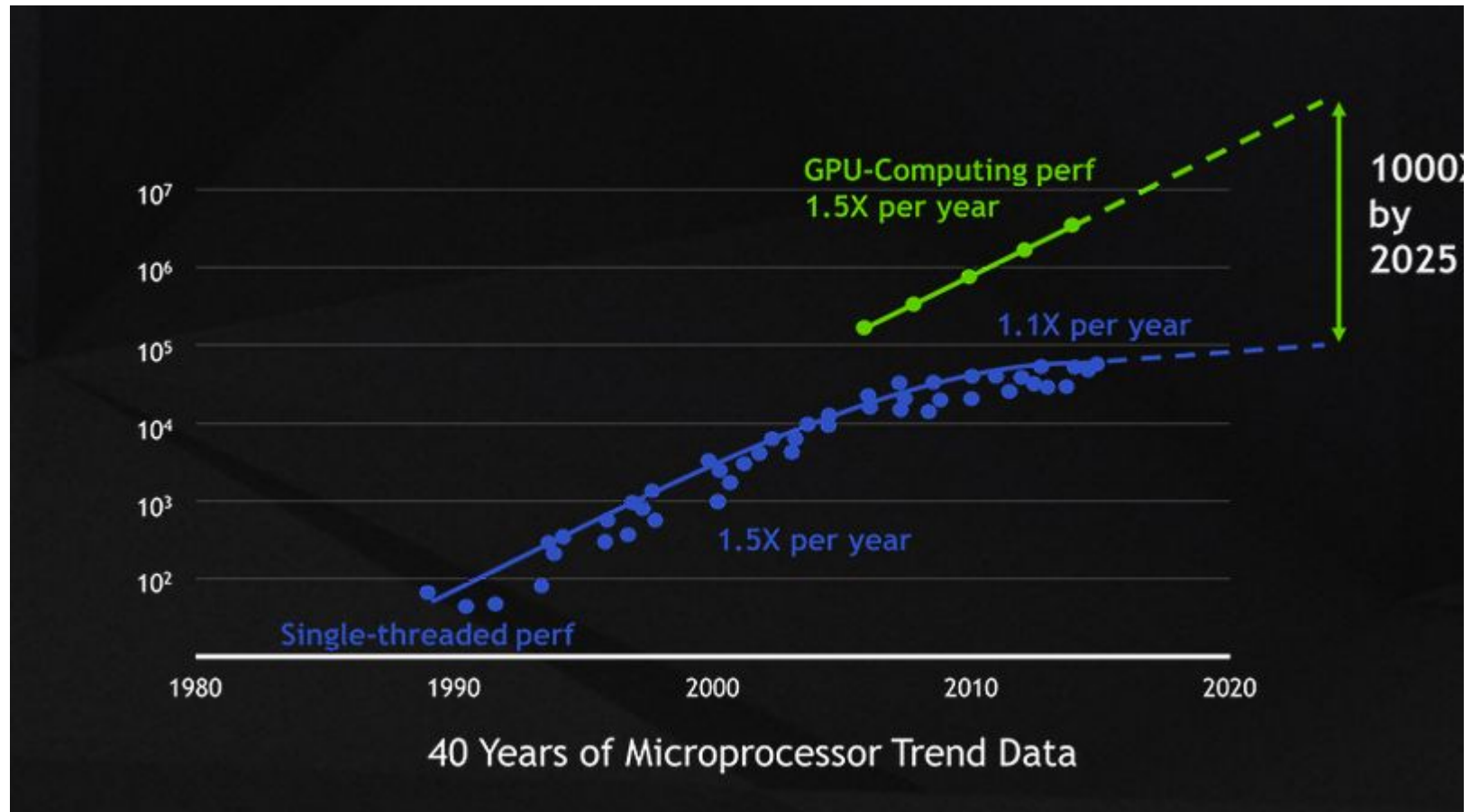


High Performance Computing on GPUs

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Why GPUs?



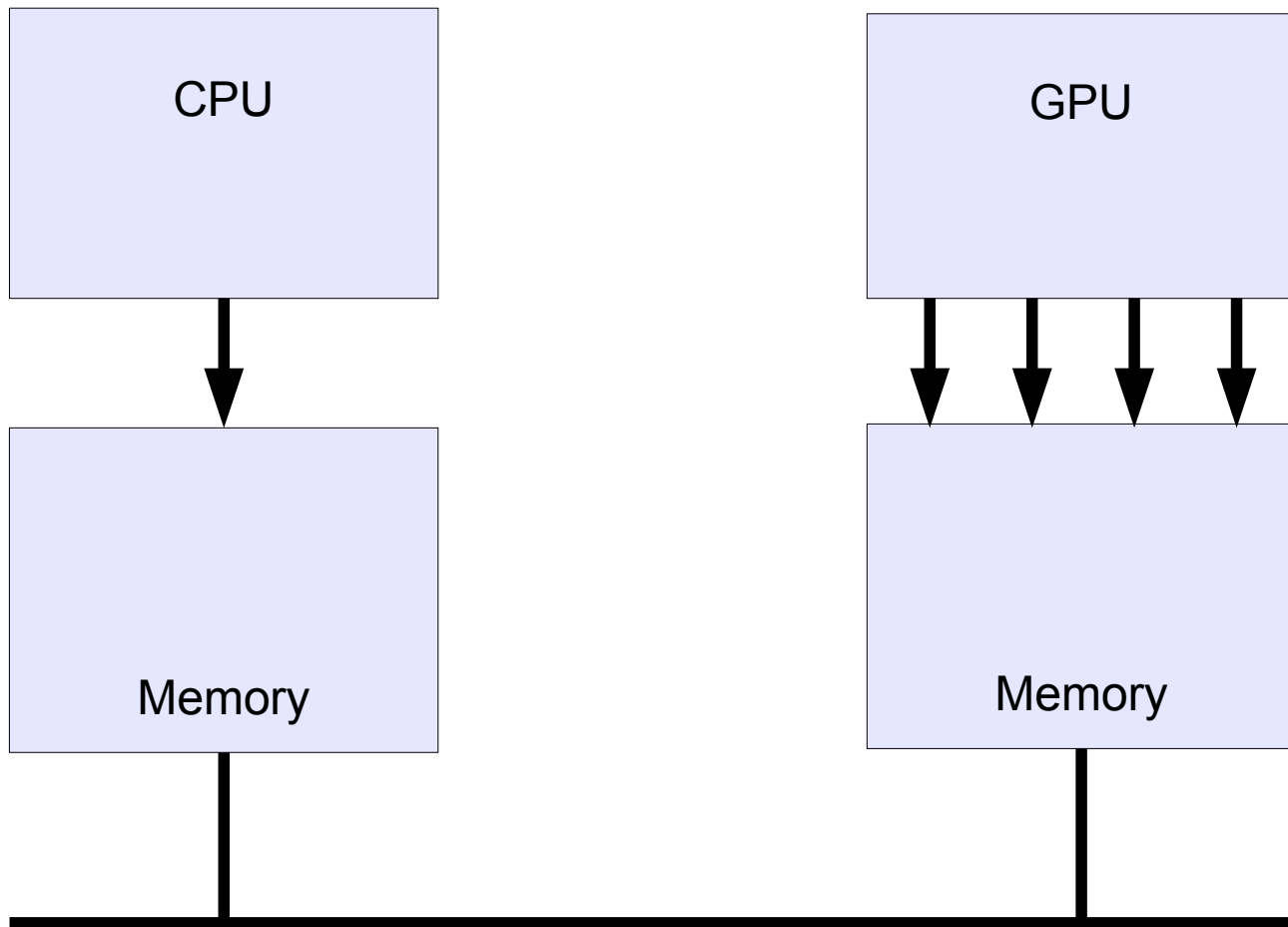
(from NVIDIA)

Is it a miracle? NO!

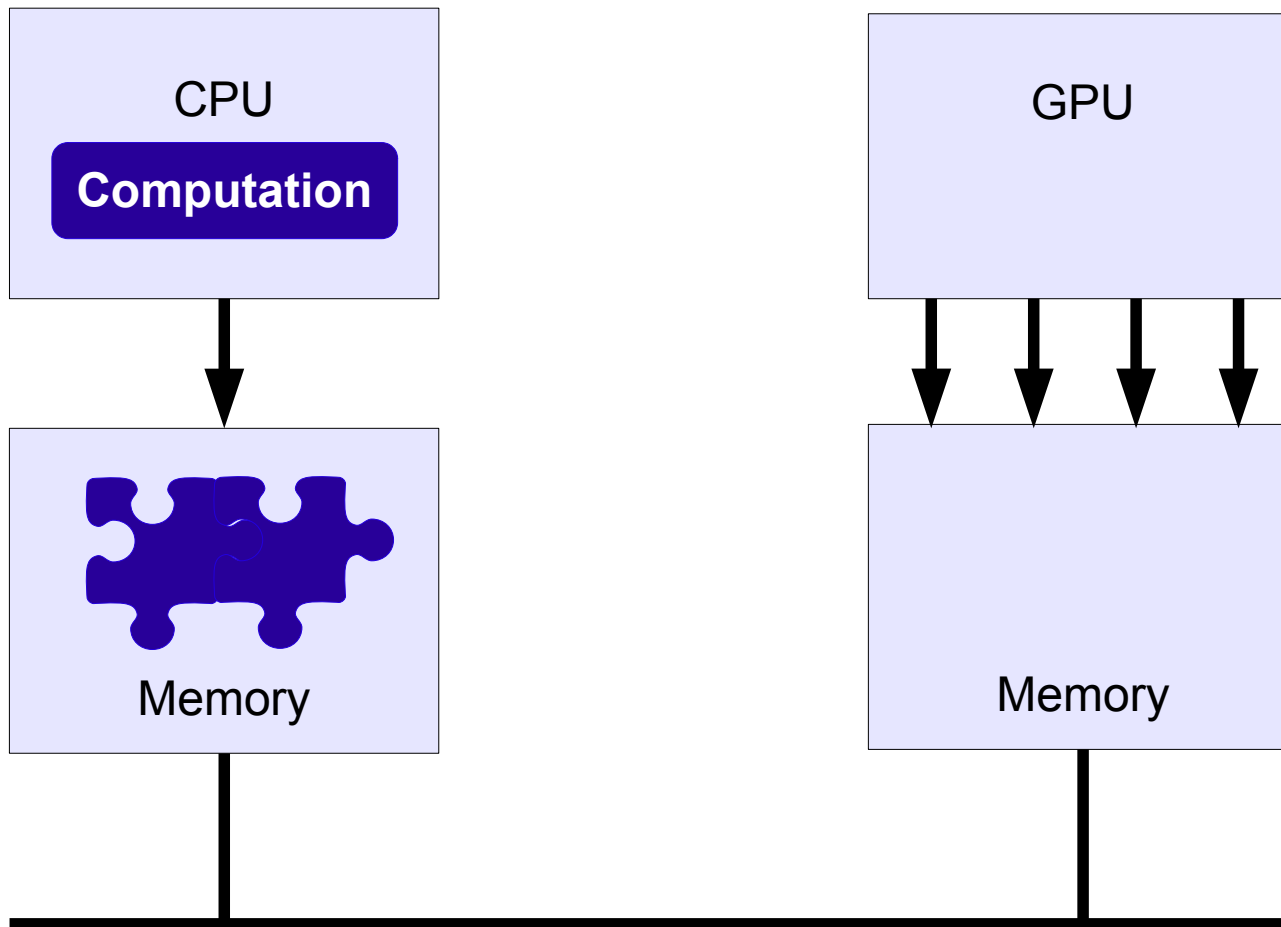
- Architectural solutions prefers parallelism!
- Example problem – I have 100 apples to eat
 - 1) “high performance”: finish one apple faster
 - 2) “high throughput”: finish **all apples** faster
- The 1st option is **unsustainable**
- Performance = parallel hardware + scalable parallel program!

Simplified GPU model

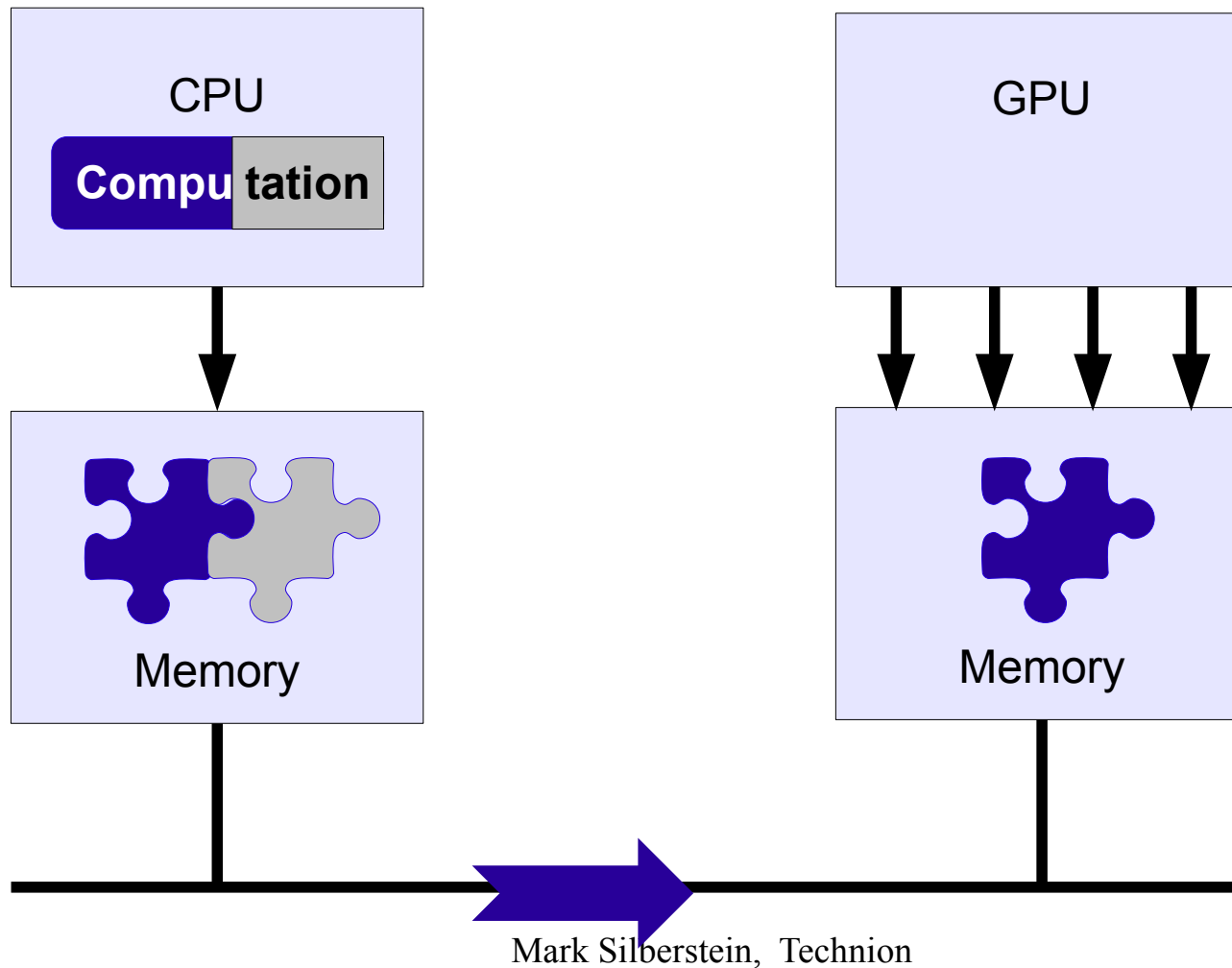
GPU 101



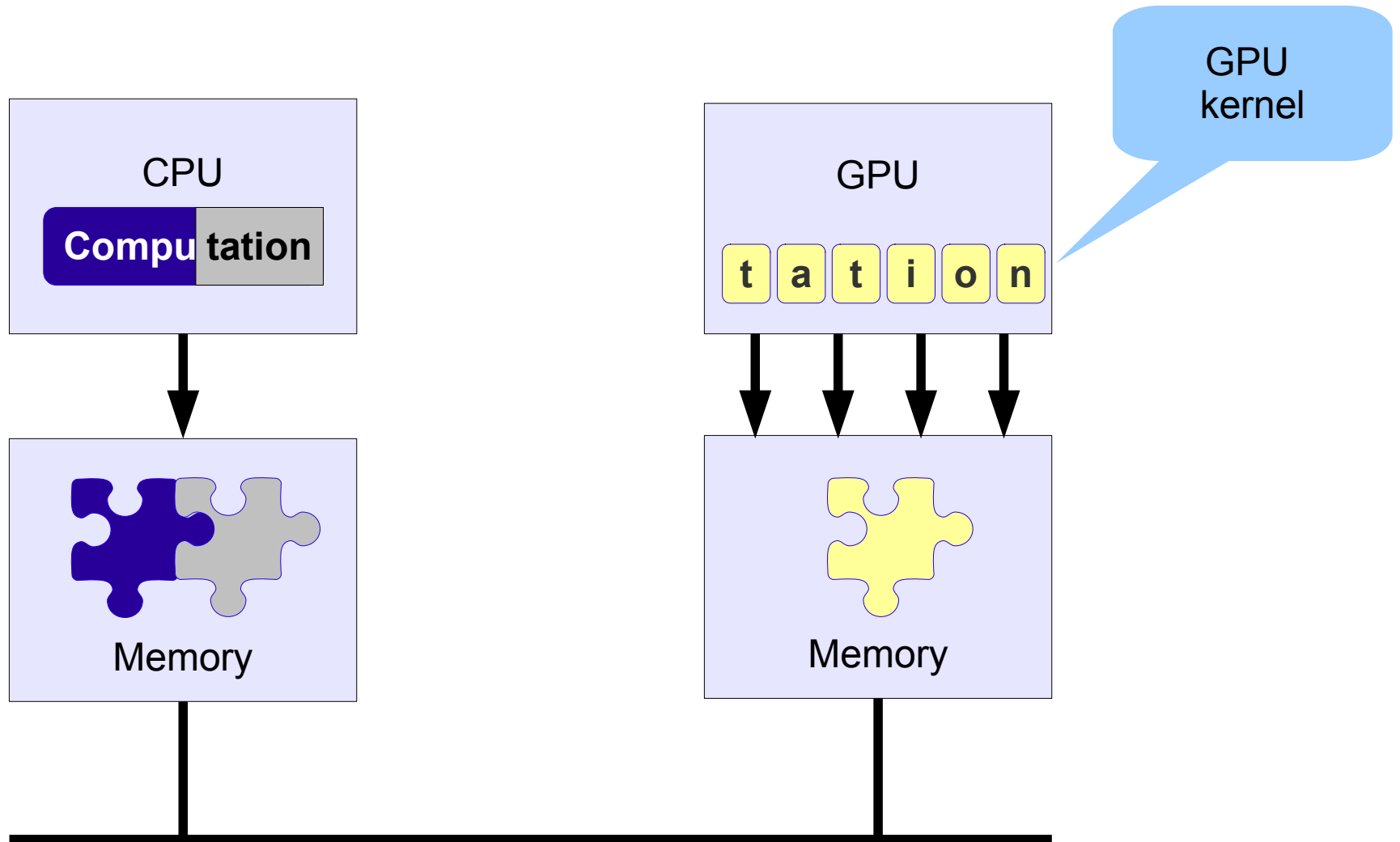
GPU is a co-processor



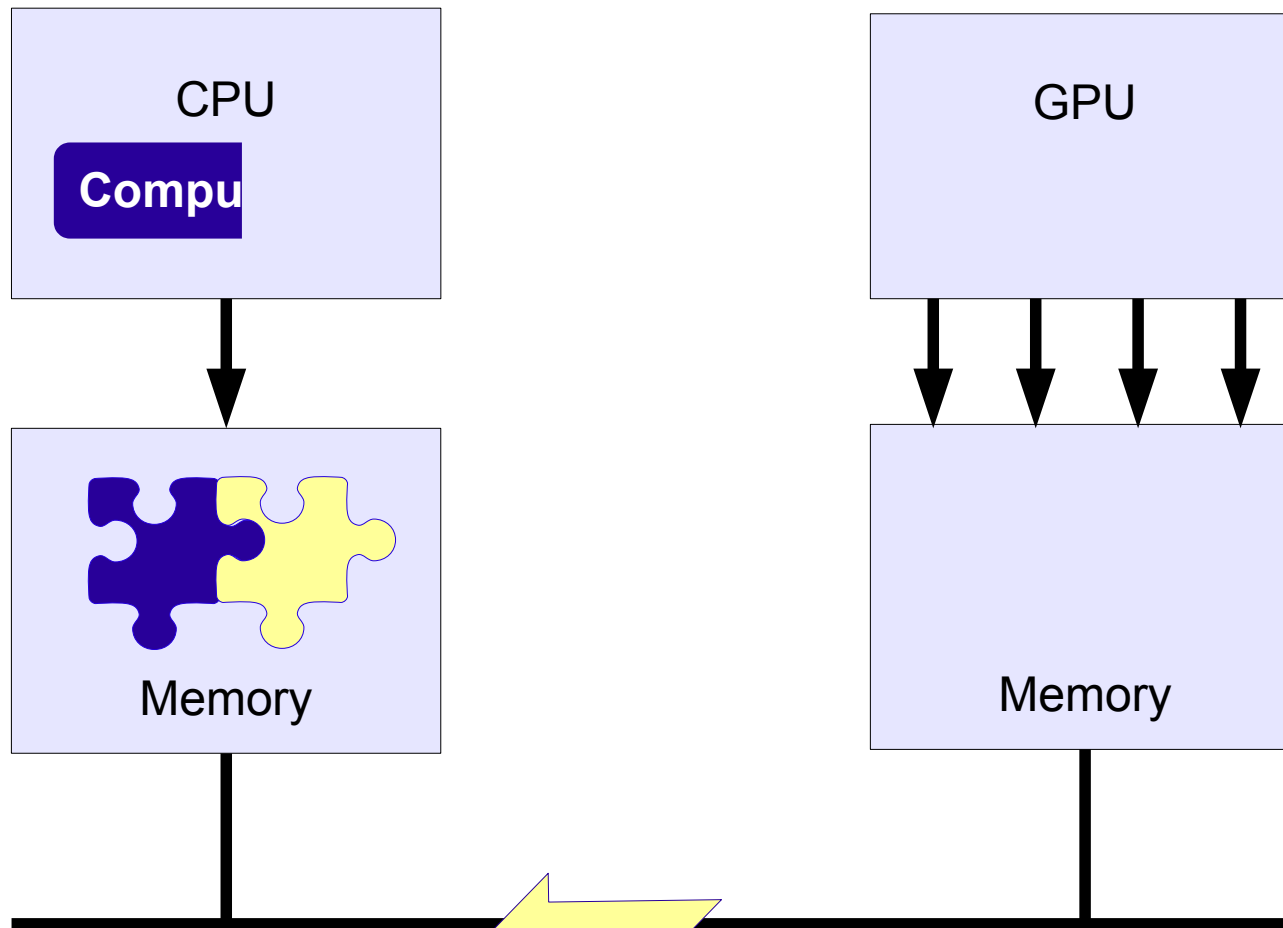
GPU is a co-processor



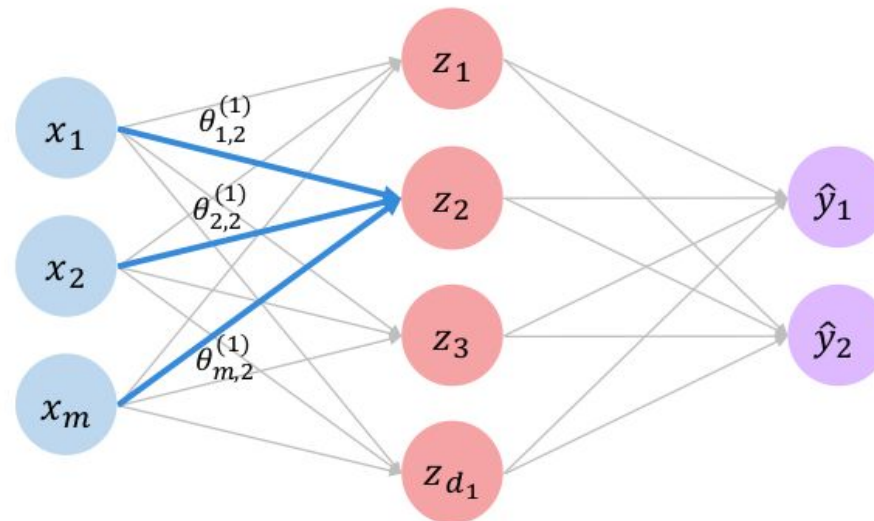
GPU is a co-processor



GPU is a co-processor



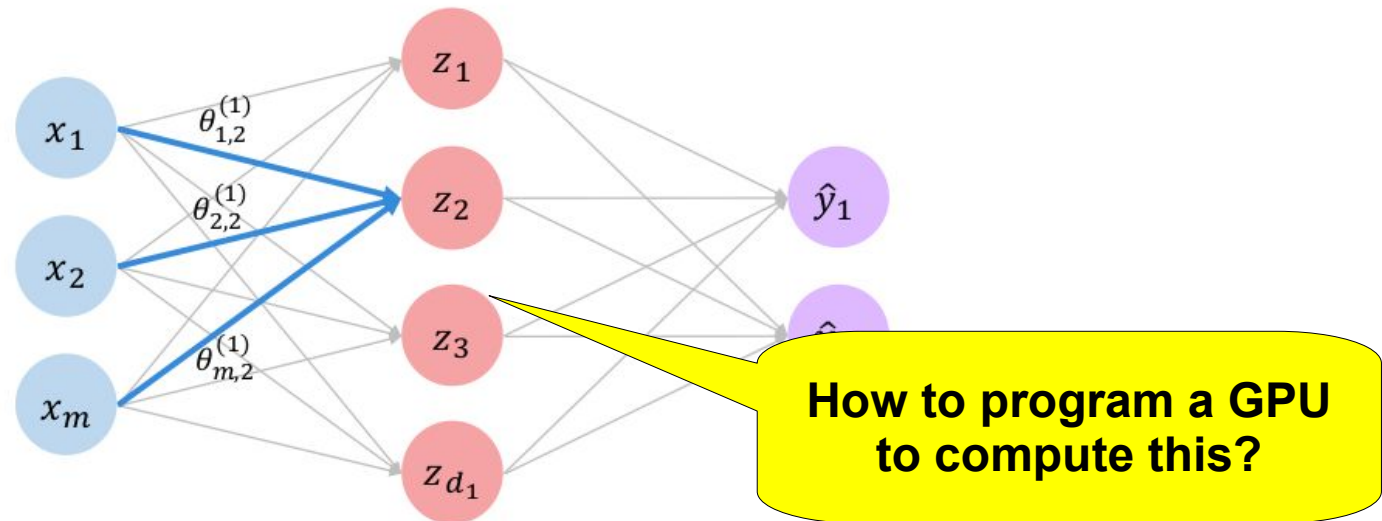
GPUs in ML – Linear Algebra Accelerators



$$\begin{aligned} z_2 &= \theta_{0,2}^{(1)} + \sum_{j=1}^m x_j \theta_{j,2}^{(1)} \\ &= \theta_{0,2}^{(1)} + x_1 \theta_{1,2}^{(1)} + x_2 \theta_{2,2}^{(1)} + x_m \theta_{m,2}^{(1)} \end{aligned}$$

Compute Matrix Product
ON GPU

GPUs in ML – Linear Algebra Accelerators



$$\begin{aligned} z_2 &= \theta_{0,2}^{(1)} + \sum_{j=1}^m x_j \theta_{j,2}^{(1)} \\ &= \theta_{0,2}^{(1)} + x_1 \theta_{1,2}^{(1)} + x_2 \theta_{2,2}^{(1)} + x_m \theta_{m,2}^{(1)} \end{aligned}$$

Simple GPU program: exploiting data parallelism

- Idea: **same set** of operations is applied to different data chunks *in parallel*
- Algorithmic challenge – identify data-parallel tasks
- Implementation
 - Every *thread* runs the same code on different data chunks.
 - GPU concurrently runs thousands of parallel threads

Vector sum $C=A+B$

- Sequential algorithm

For every element i

$$C[i]=A[i]+B[i]$$

Vector sum $C=A+B$

- Sequential algorithm

For every i

$$C[i]=A[i]+B[i]$$

- Parallel algorithm

In parallel For every i

$$C[i]=A[i]+B[i]$$

Implementation for a vector of length 1024

- **GPU kernel** (this program runs in every thread)

$C[\text{threadId}] = A[\text{threadId}] + B[\text{threadId}]$



Per-thread hardware-supplied ID

Implementation for a vector of length 1024

- GPU kernel

$C[\text{threadId}] = A[\text{threadId}] + B[\text{threadId}]$

- CPU

1. Allocate three arrays (in GPU memory)

2. Make data accessible to GPU (CPU->GPU copy)

3. Invoke kernel with 1024 threads

4. Wait until complete and make data accessible to CPU (GPU->CPU copy)

Complete example

CPU:

```
void vector_sum(float* A, float* B, float* C, int n)
{
    float* gA=GPU_get_reference(A);
    float* gB=GPU_get_reference(B);
    float* gC=GPU_allocate_mem(n);

    GPU_set_num_threads(n);
    // GPU will invoke n threads
    GPU_run(vector_sum_kernel(gA,gB,gC));
    GPU_retrieve(C,gC);
}
```

GPU:

```
void vector_sum_kernel(float* gA, float* gB, float*gC)
{
    int my=HardwareThreadID;
    gC[my]=gA[my]+gB[my];
}
```

Complete example

CPU:

```
void vector_sum(float* A, float* B, float* C, int n)
{
    float* gA=GPU_get_reference(A);
    float* gB=GPU_get_reference(B);
    float* gC=GPU_allocate_mem(n);

    GPU_set_num_threads(n);
    // GPU will invoke n threads
    GPU_run(vector_sum_kernel(gA,gB,gC));
    GPU_retrieve(C,gC);
}
```

GPU integration

GPU:

```
void vector_sum_kernel(float* gA, float*
{
    int my=HardwareThreadID;
    gC[my]=gA[my]+gC[my];
}
```

GPU programming

Complete example

```
CPU:  
void vector_sum(float* A, float* B  
{  
    float* gA=GPU_get_reference(A)  
    float* gB=GPU_get_reference(B);  
    float* gC=GPU_allocate_mem(n);
```

Many threads

```
GPU_set_num_threads(n);  
    // GPU will invoke n threads  
GPU_run(vector_sum_kernel(gA,gB,gC));  
GPU_retrieve(C,gC);  
}
```

**Fine-grained parallelism
(1 op per thread)**

```
GPU:  
void vector_sum_kernel(float*  
{  
    int my=HardwareThreadID;  
    gC[my]=gA[my]+gB[my];  
}
```

BUT!

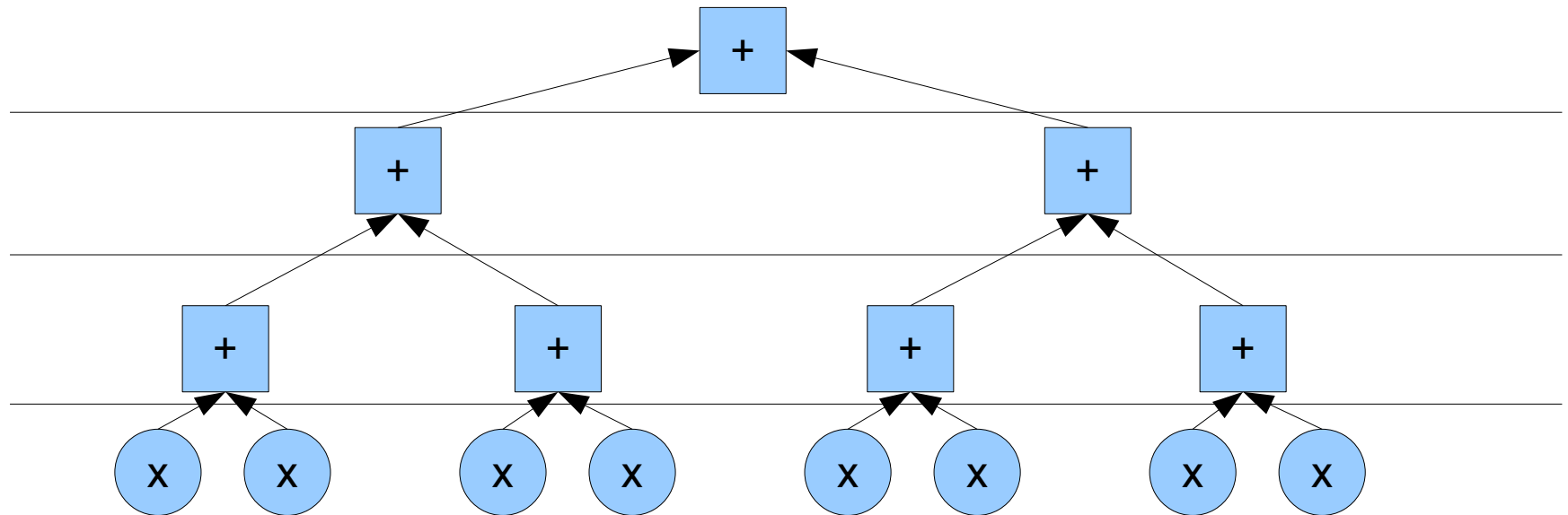
- Vector sum is simple – purely data parallel
- What if we need coordination between tasks

Example: parallel dot product

BUT!

- Vector sum is simple – purely data parallel
- What if we need coordination between tasks

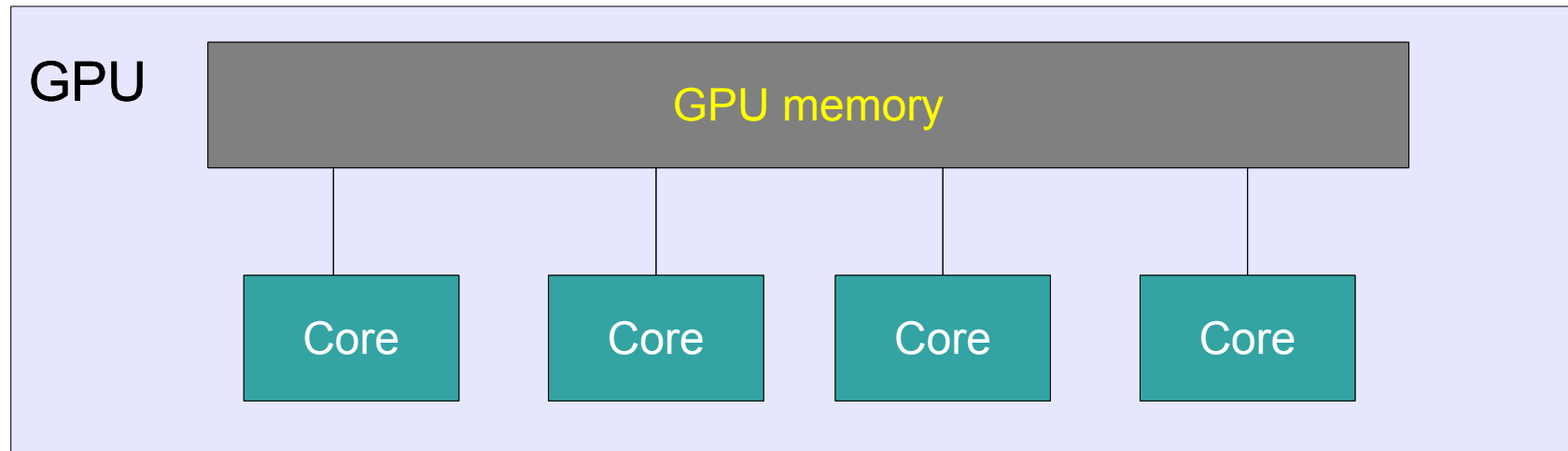
Example: parallel dot product



GPU hardware

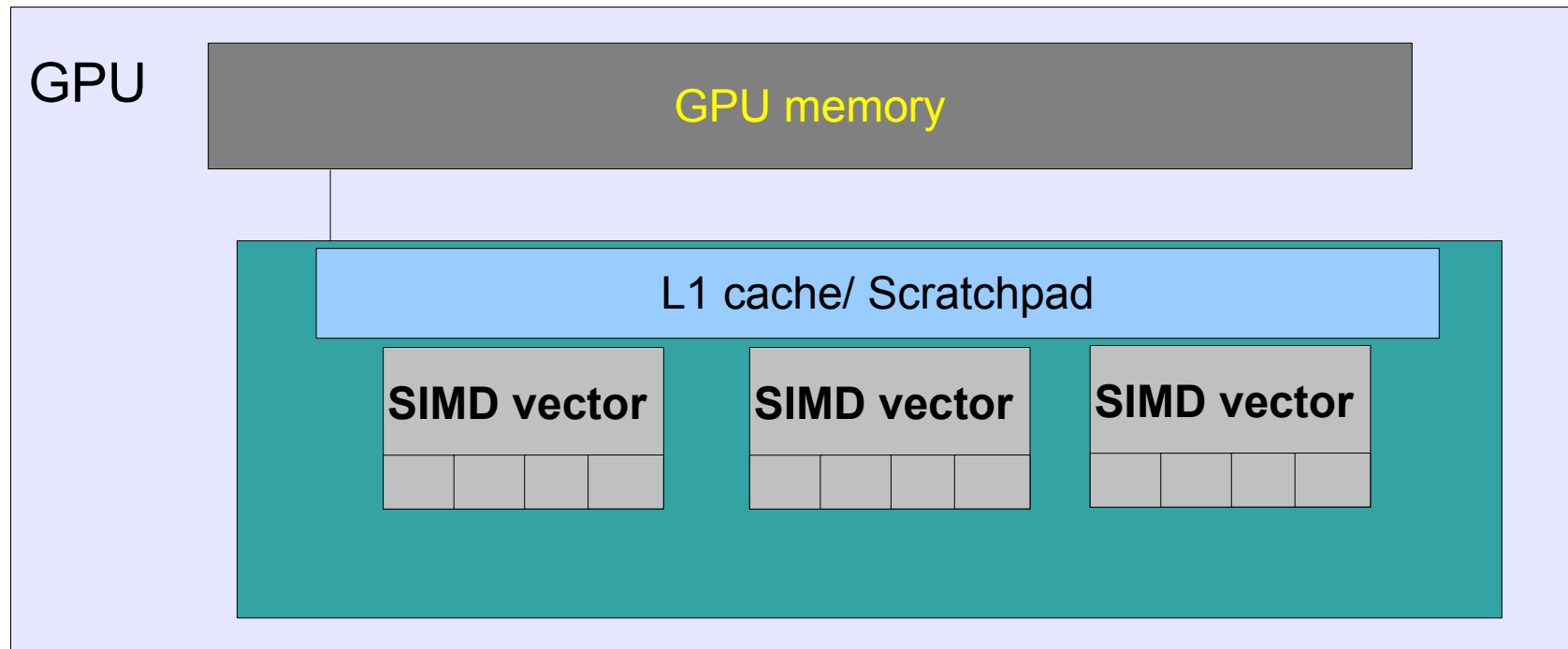
GPU hardware parallelism

1. Multi-core



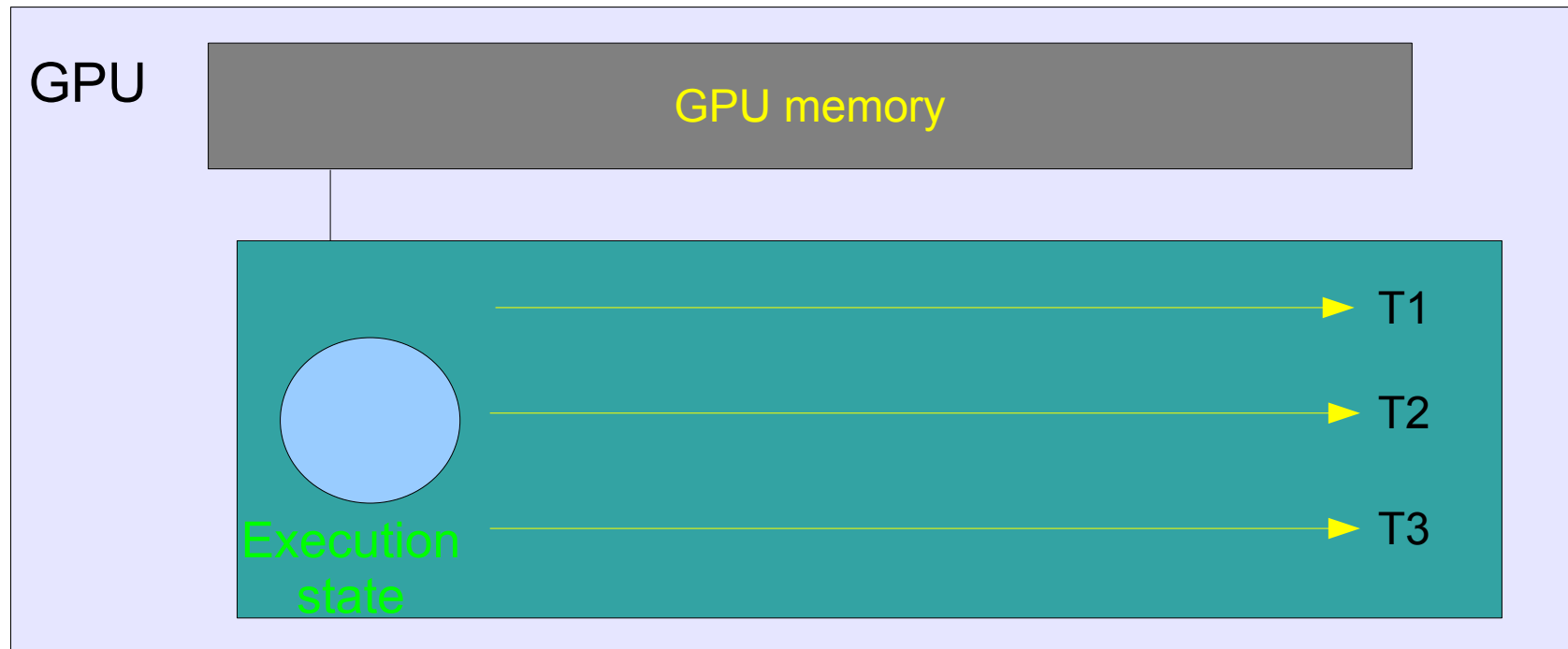
GPU hardware parallelism

2. SIMD



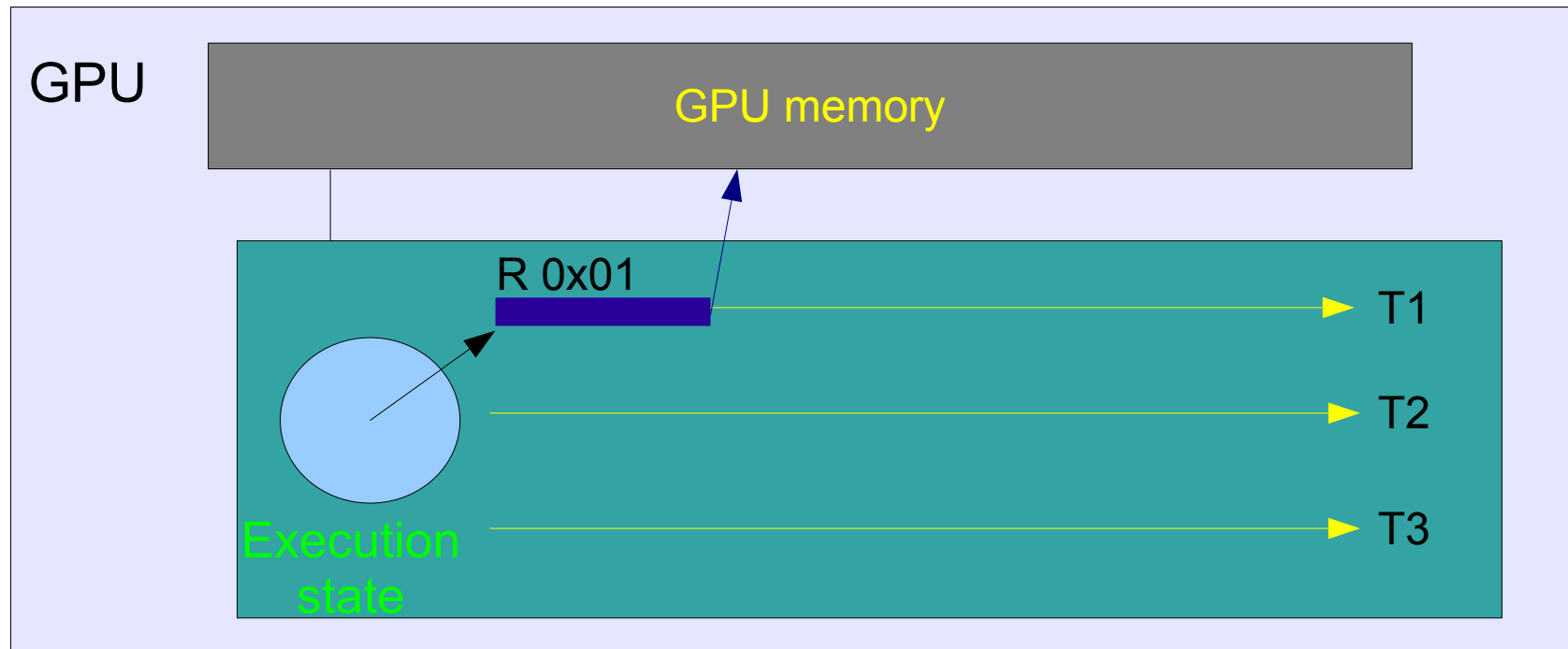
GPU hardware parallelism

3. Hardware multithreading



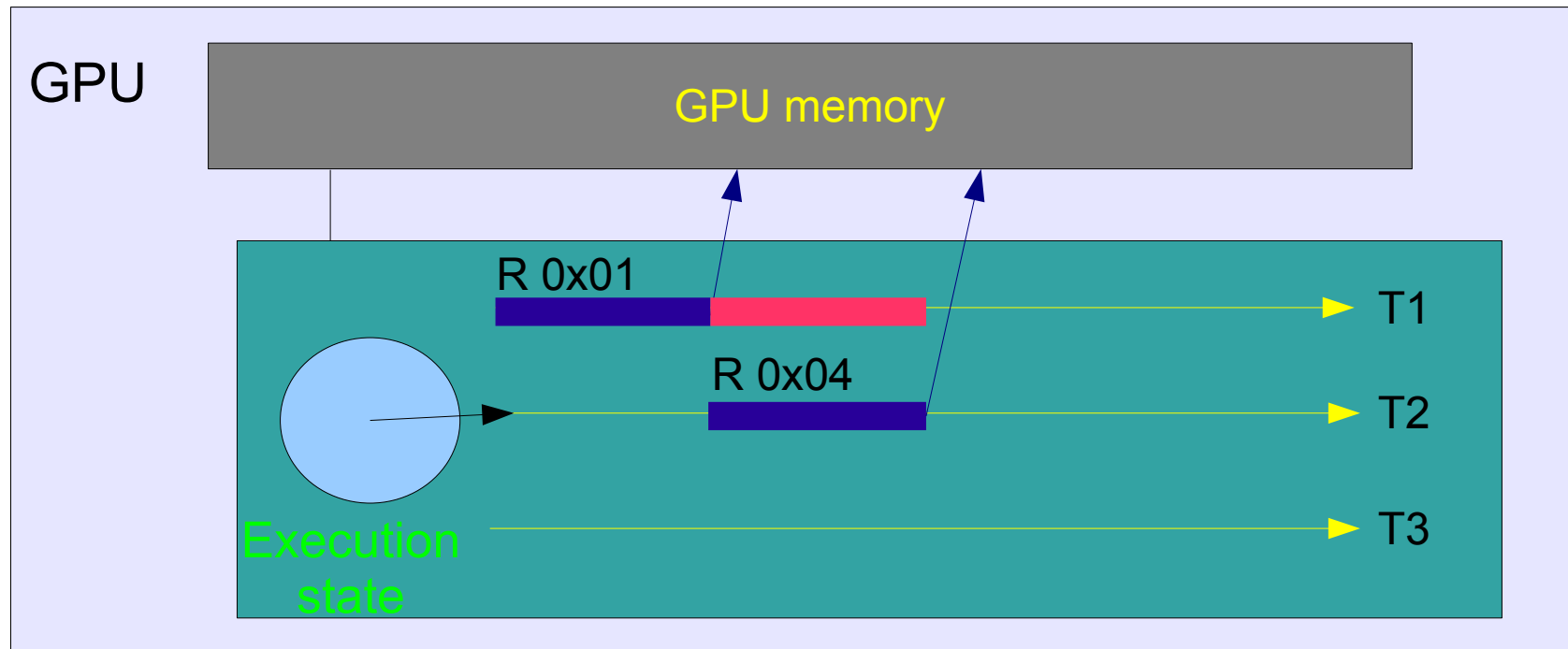
GPU hardware parallelism

3. Hardware multithreading



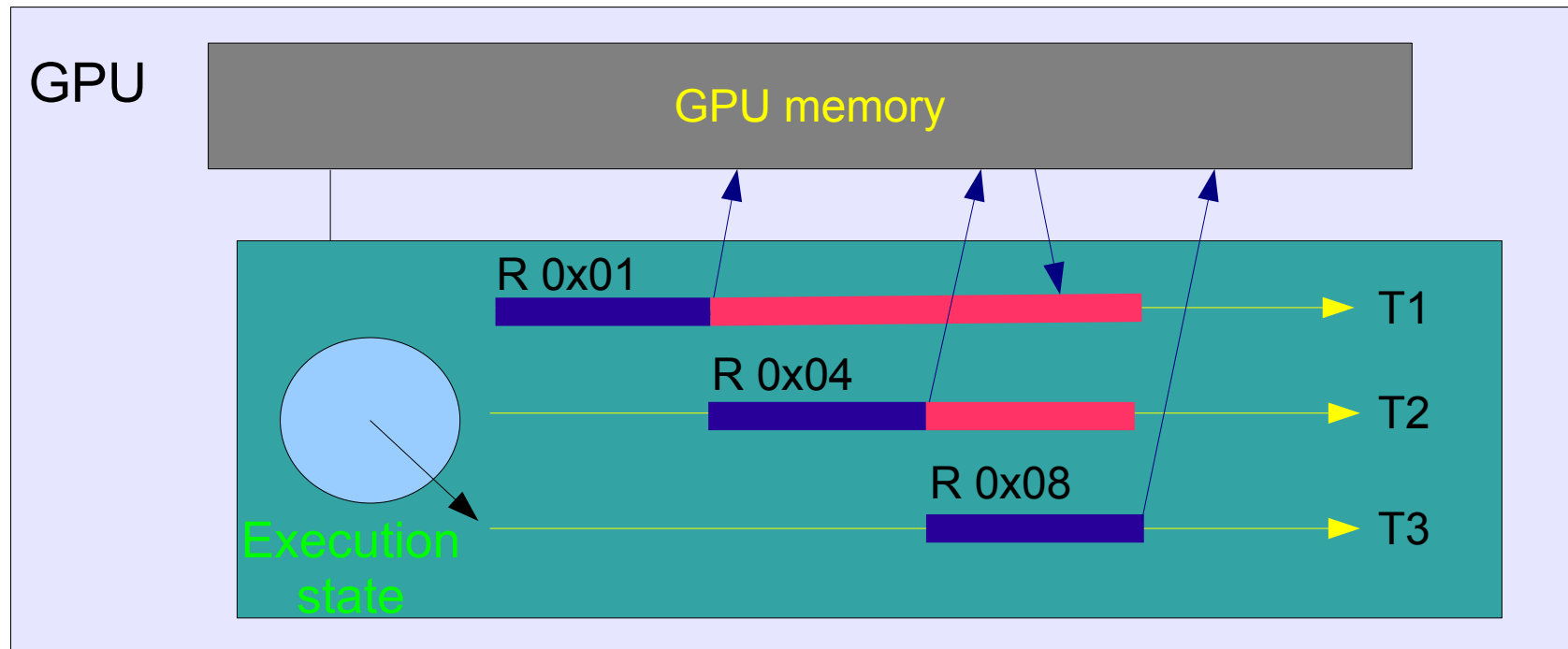
GPU hardware parallelism

3. Hardware multithreading



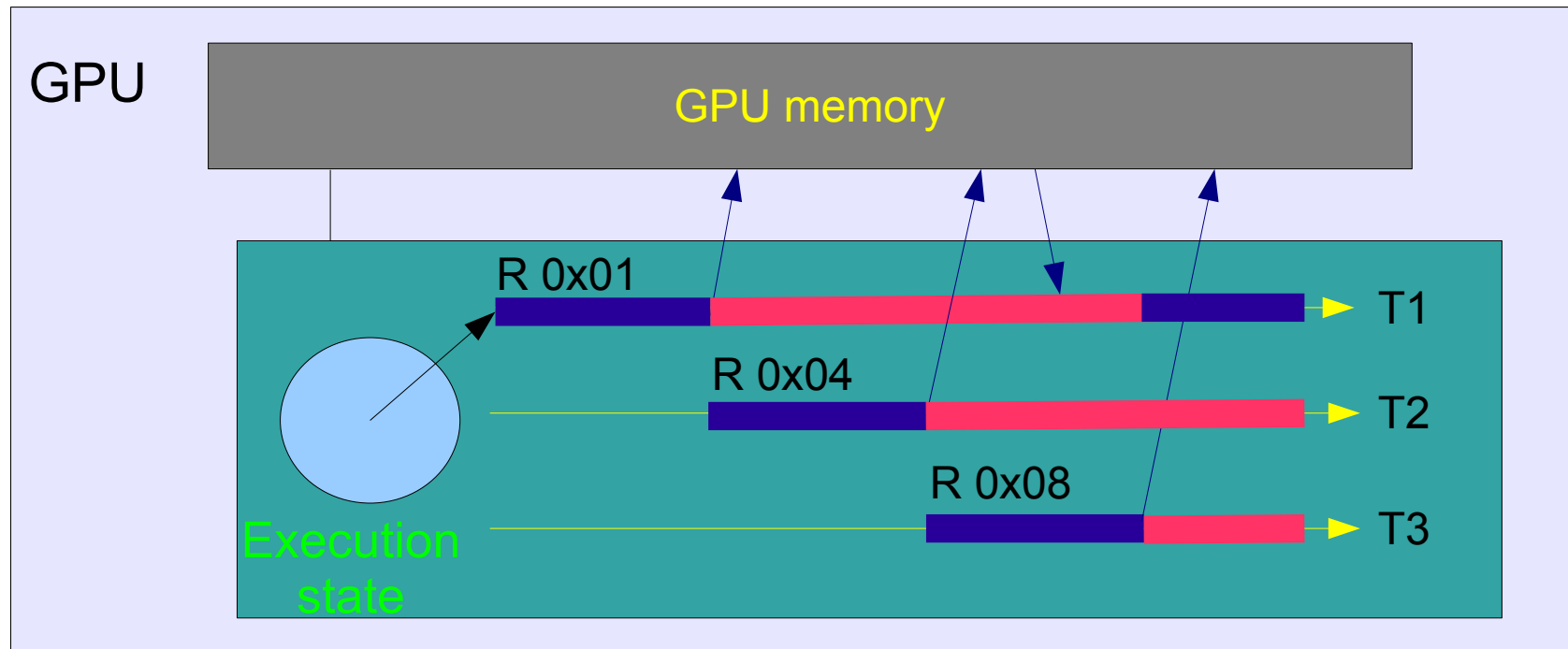
GPU hardware parallelism

3. Hardware multithreading

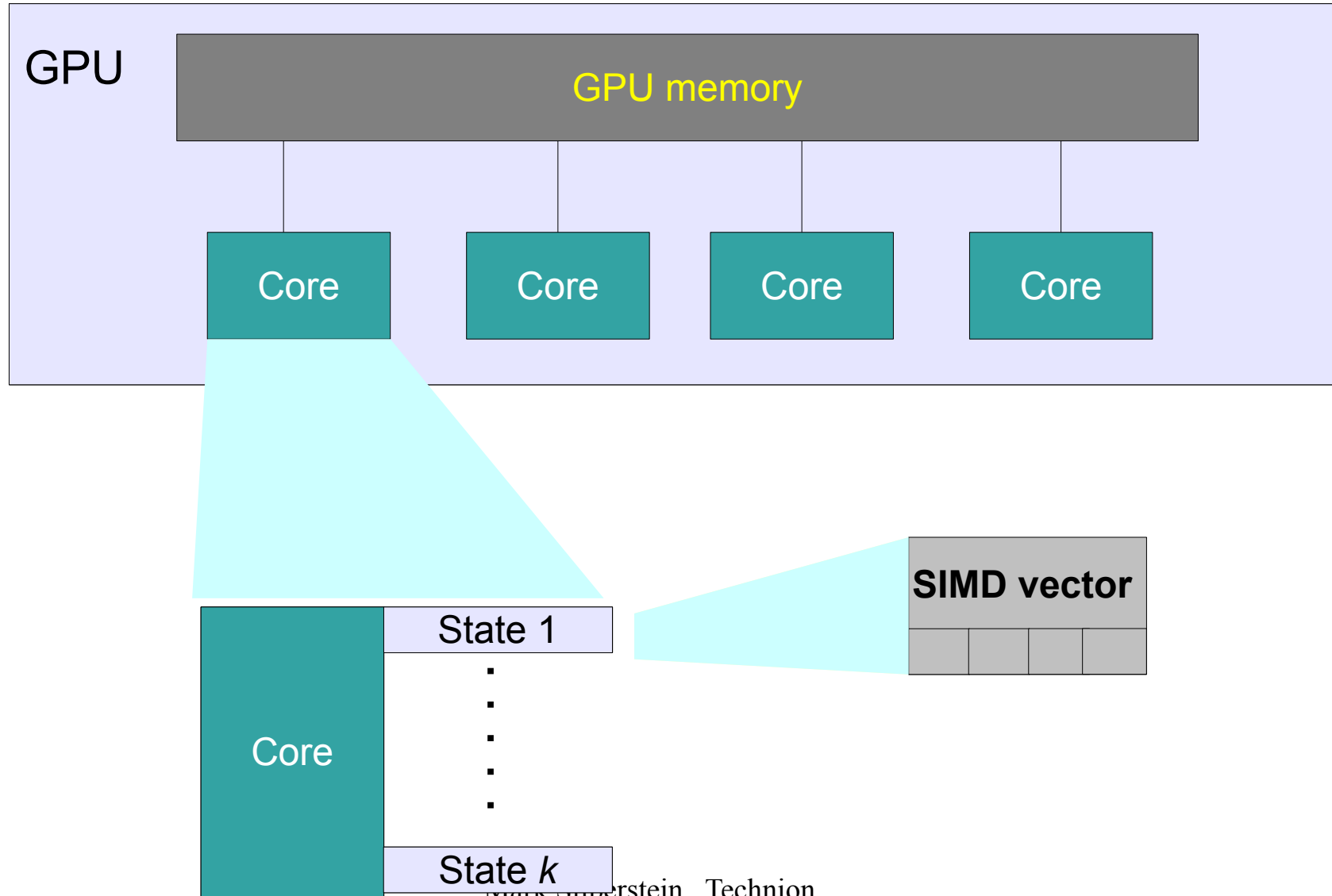


GPU hardware parallelism

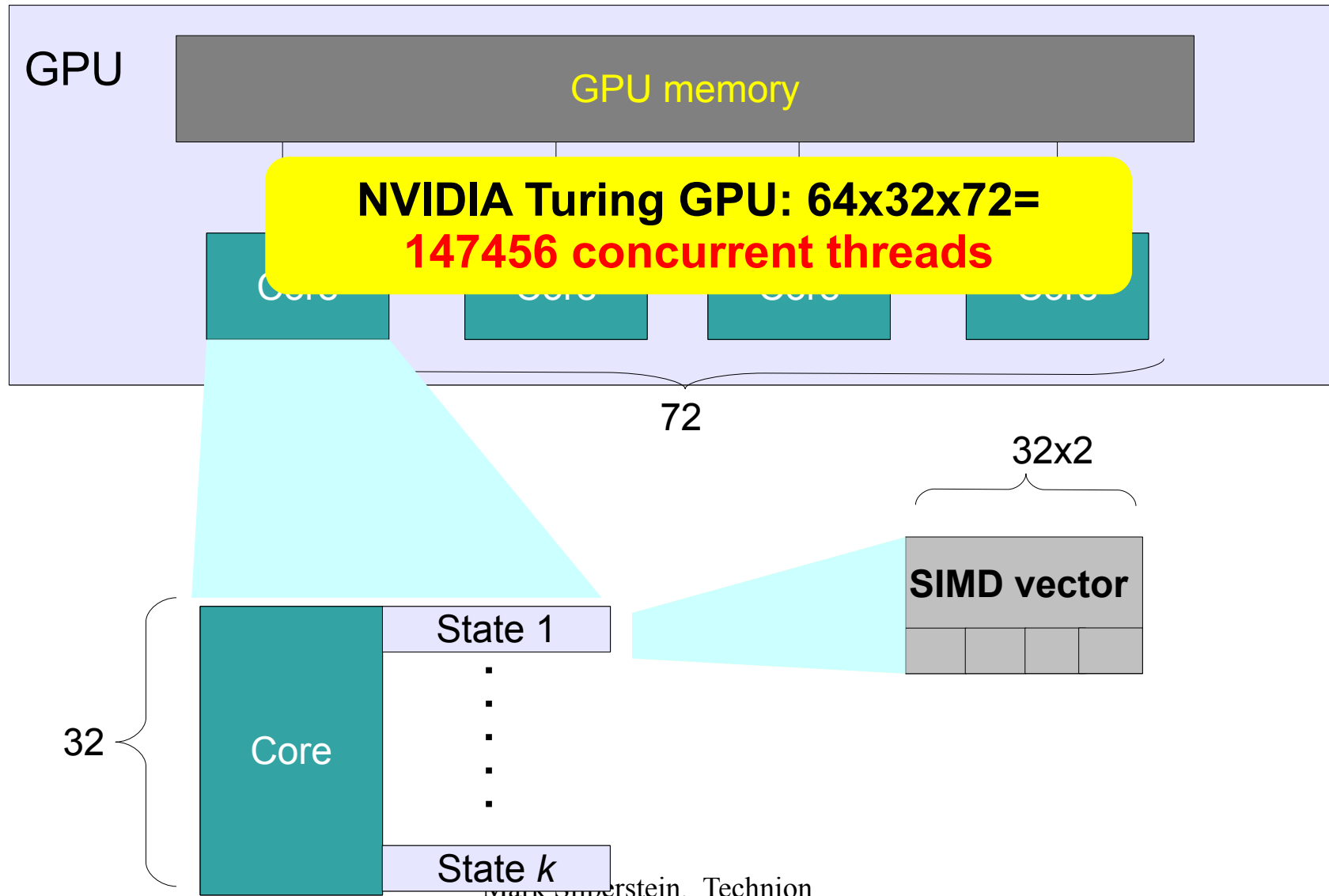
3. Hardware multithreading



Putting it all together: 3 levels of hardware parallelism



Takeaway 1: 100,000-s of concurrent threads!



Requirements for allowing fast GPU execution

- We have enough parallelism
- We have enough space to store state per thread
- We have enough bandwidth to memory
- We have efficient *scheduler* to manage threads

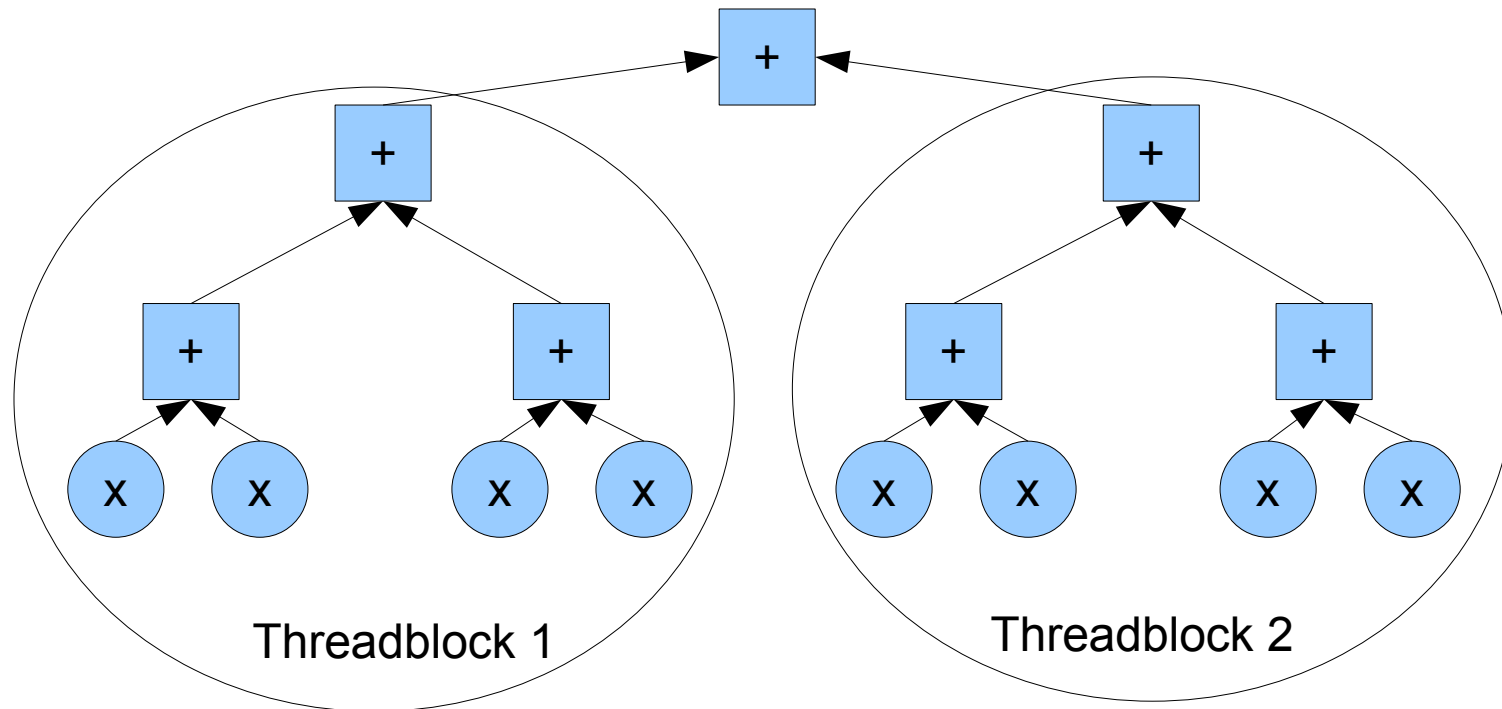
How GPU manages threads?

- Application threads are grouped into threadblocks
- Programmer defines the number of threads / threadblocks for each run

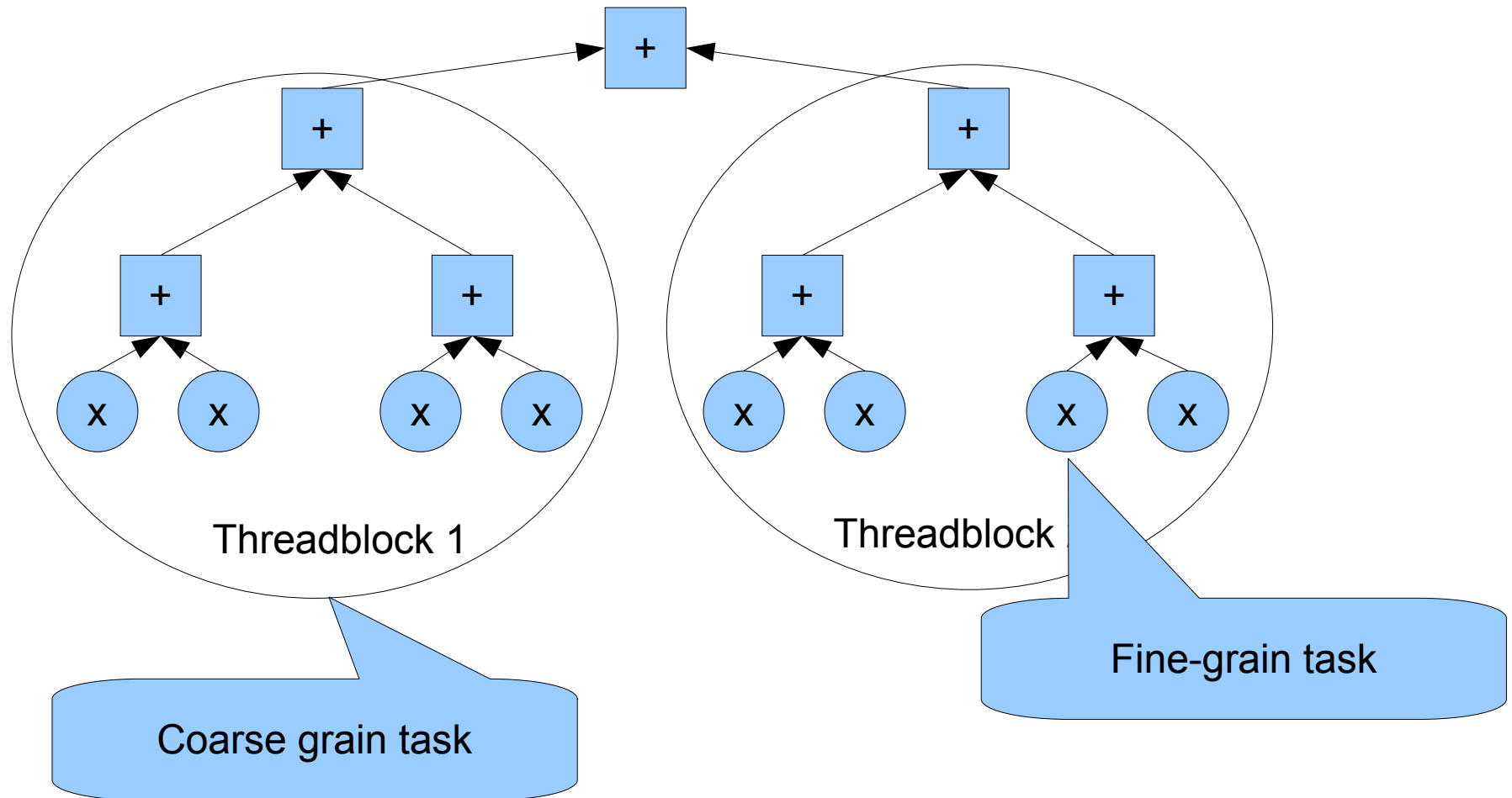
Threadblock is a building block of GPU algorithms

- Threads inside a threadblock can communicate efficiently!
 - Share a small fast scratchpad memory
 - Can be synchronized via barriers
- Threadblocks are independent
- A program consists of many threadblocks

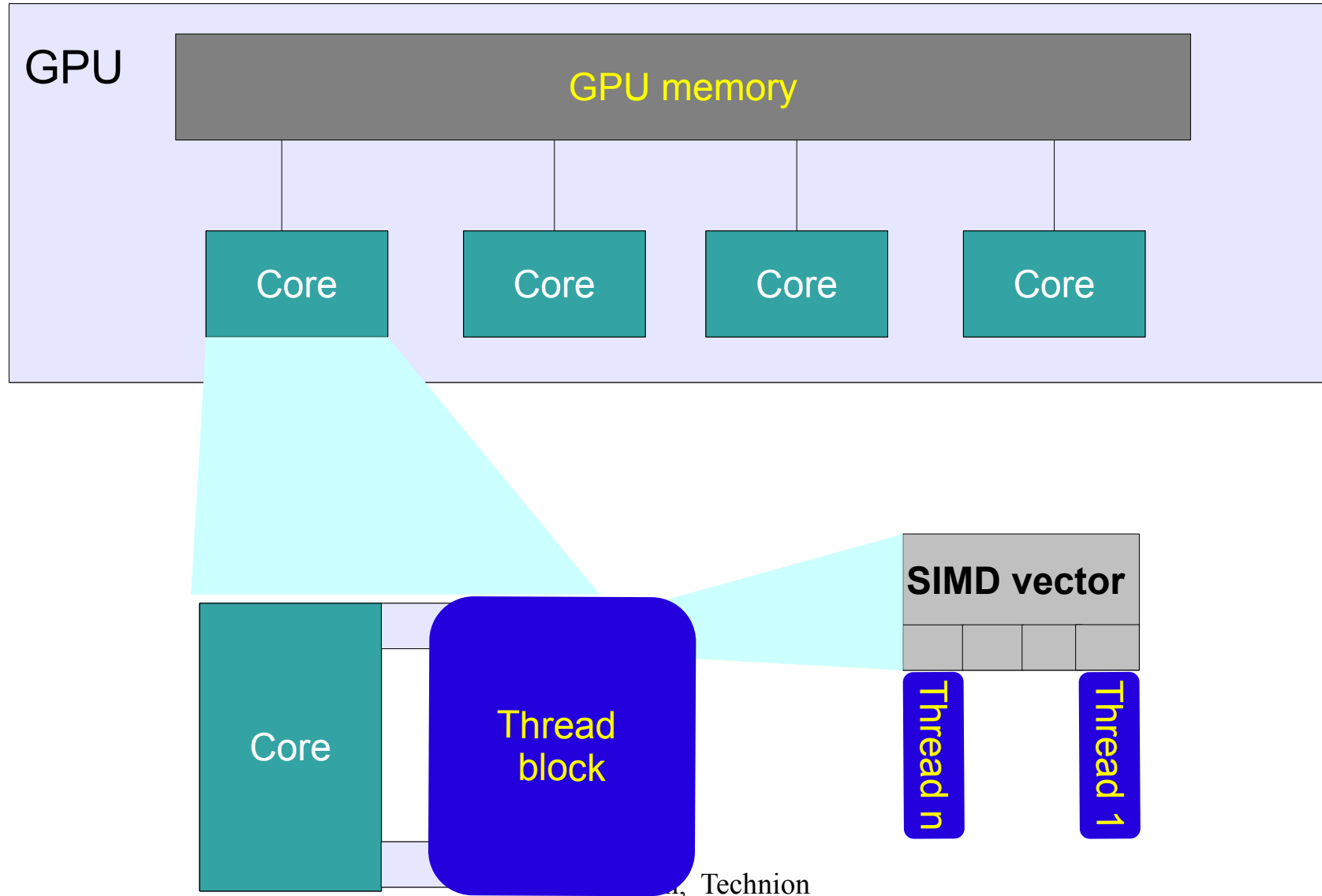
Dot-product: hierarchical parallelization Decomposing into threadblocks



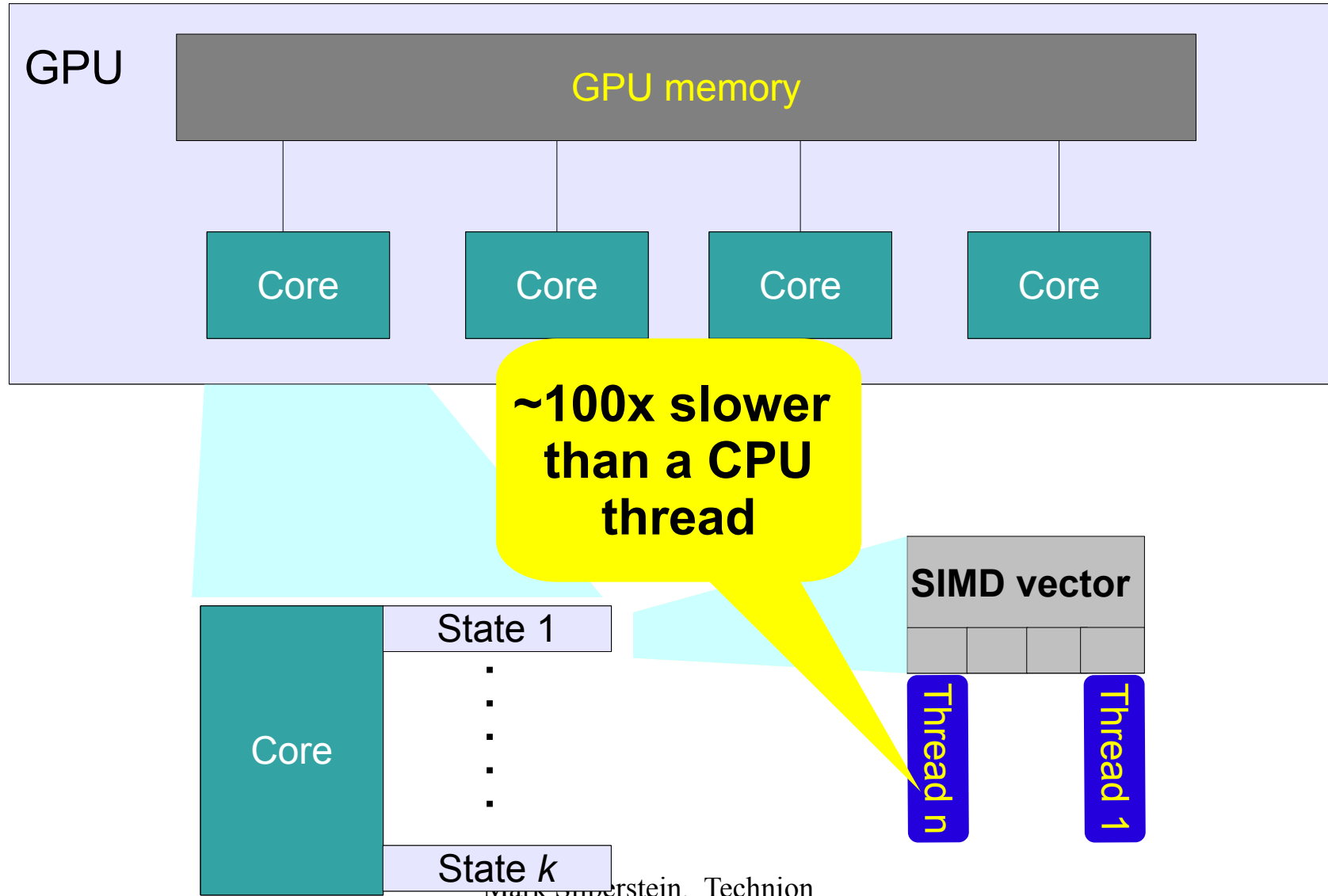
Takeaway: parallelism hierarchy



Software-Hardware mapping



Takeaway 2: One thread is slow



Dot product

```
void vector_dotproduct_kernel(float* gA, float* gB, float* gOut)
{
    _local_ float l_res[TB_SIZE]; //local core memory

    int thid=LocalThreadID;
    int tbid=ThreadBlockID;

    int offset=tbid*WG_SIZE+thid;

    l_res[tid]=gA[offset]*gB[offset];

    BARRIER(); // wait for all products

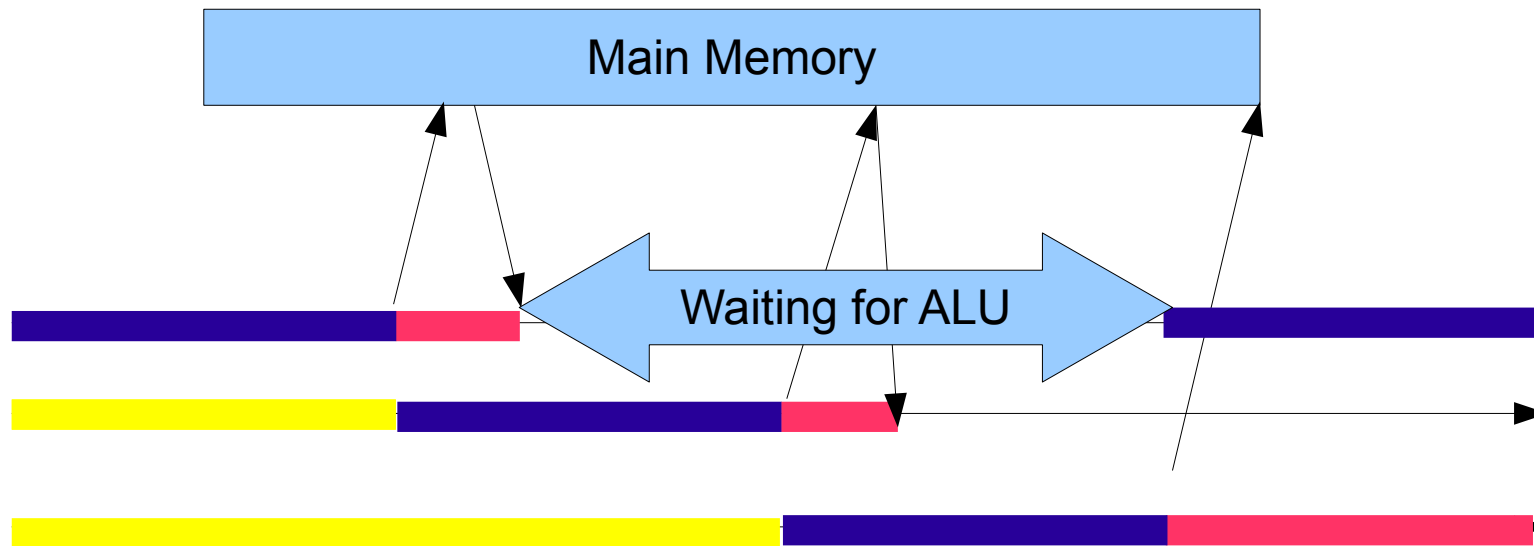
    for(int i=TB_SIZE/2; i>0;i/=2)
    {
        if (thid<i) l_res[thid]=l_res[thid]+l_res[i+thid];
        BARRIER(); // wait for all partial sums
    }
    if (thid==0) gOut[tbid]=l_res[0];
}
```

Parallelism structure of GPU programs

- Having many independent tasks is not enough
- Parallel structure should map well on hardware hierarchy

Estimating application performance on high-throughput processors

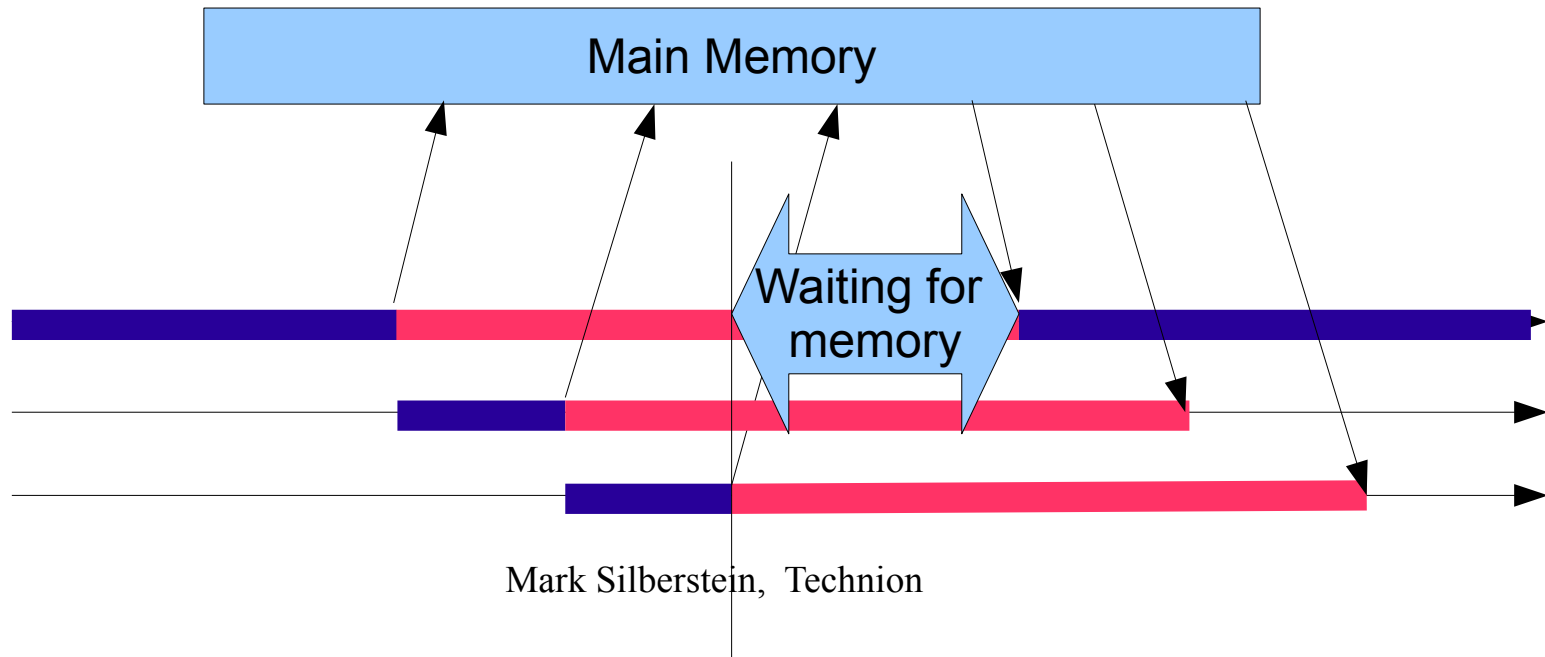
Compute-bound tasks



Performance bounded by maximum
ALU capacity

Memory-bound tasks

- We can fully utilize a processor only if data is available in an ALU on time
- How fast can the data be made available?
 - Assume infinite number of threads, ideal parallelization



Measure of ALU/memory ratio: Arithmetic intensity

- Number of OPs per **memory(**)** access
 - Vector sum: 1 operation per 3 accesses. $A=1/3$

Upper bound on performance

- For memory bound algorithms only
(why?)

$$\text{Perf} = \text{GPUMemBandwidth} * A$$

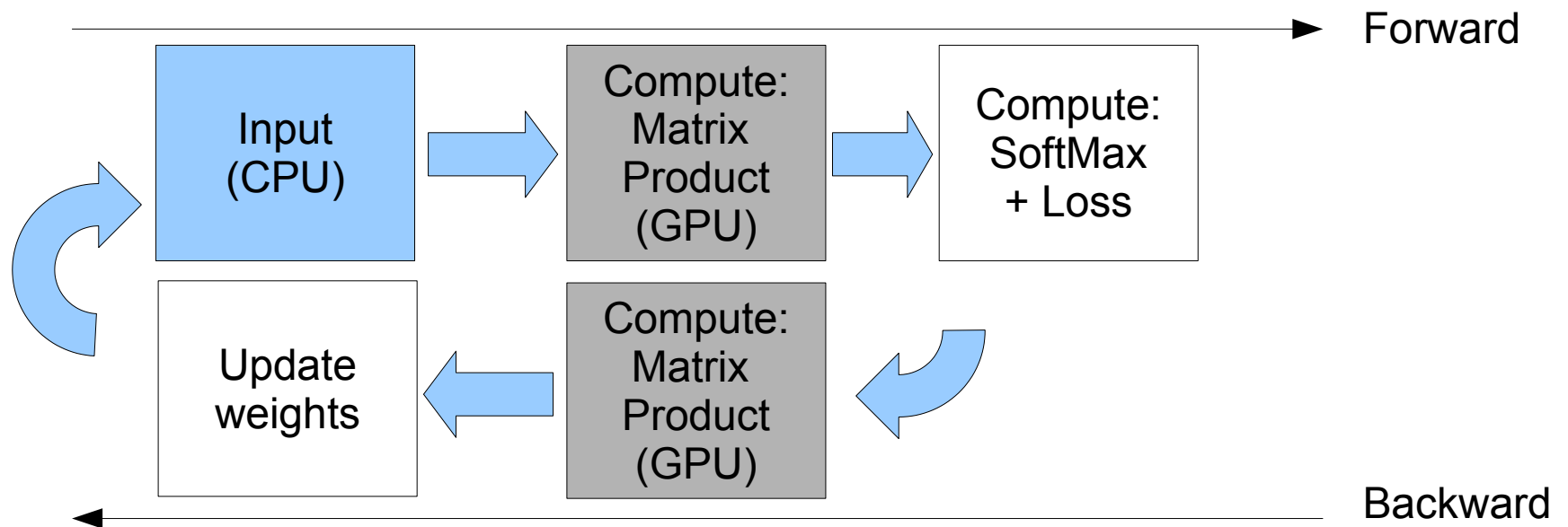
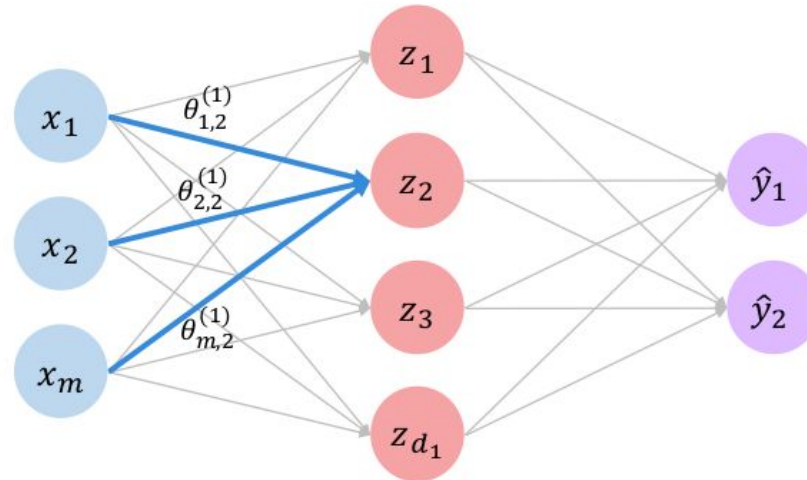
$$\text{Bytes/sec} * \text{Ops/Bytes} = \text{Ops/sec}$$

Vector sum performance estimate

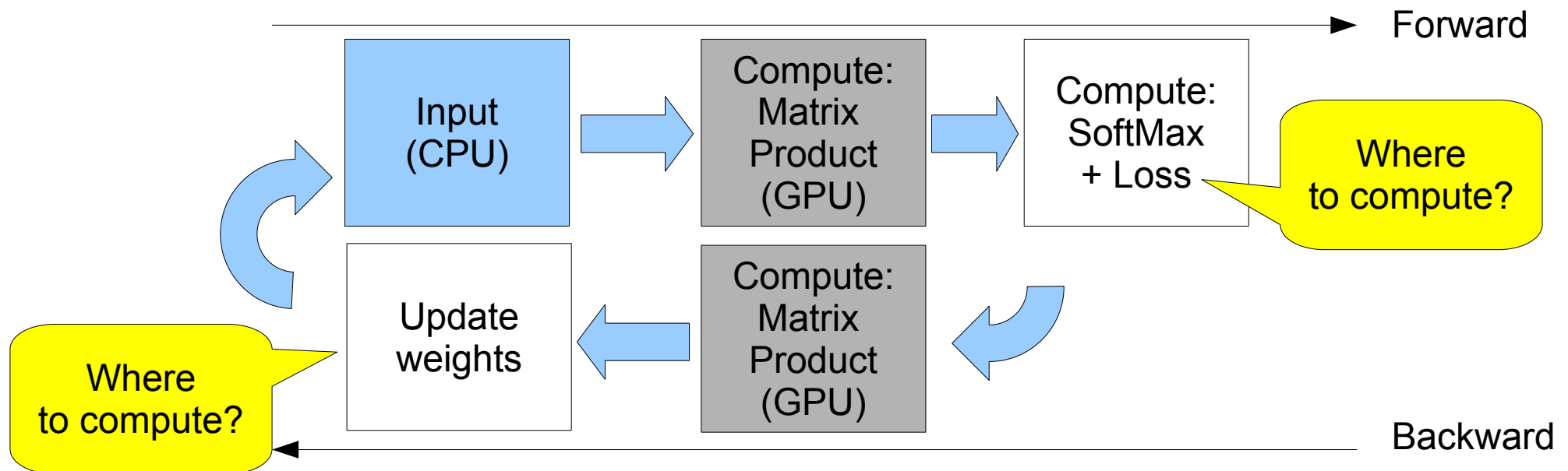
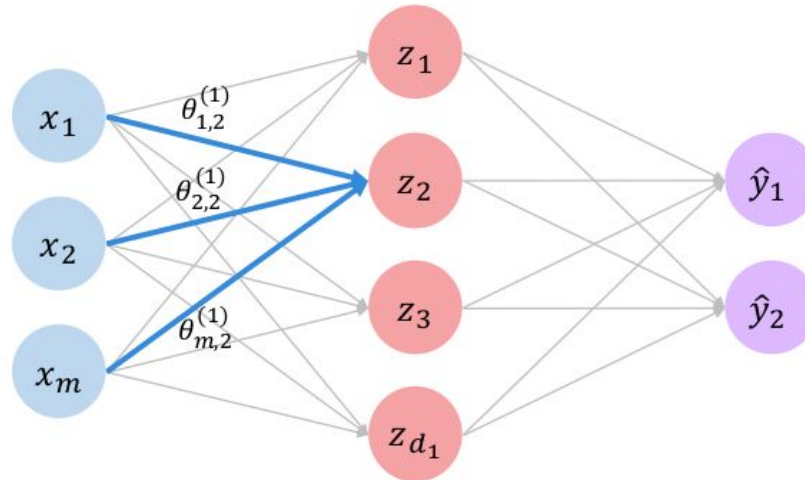
- Example: sum of vectors, GPU GTX Titan
- $A=1/(3*(4 \text{ bytes}))$, MemBW= $\sim 70 \text{ GFloat/s}$:
Performance= $\sim 23 \text{ GFLOP/s}$
- For comparison: raw capacity: 4.5 TFLOPs
- **Only 0.5% of computing capacity utilized!!**

Integrating GPUs with applications

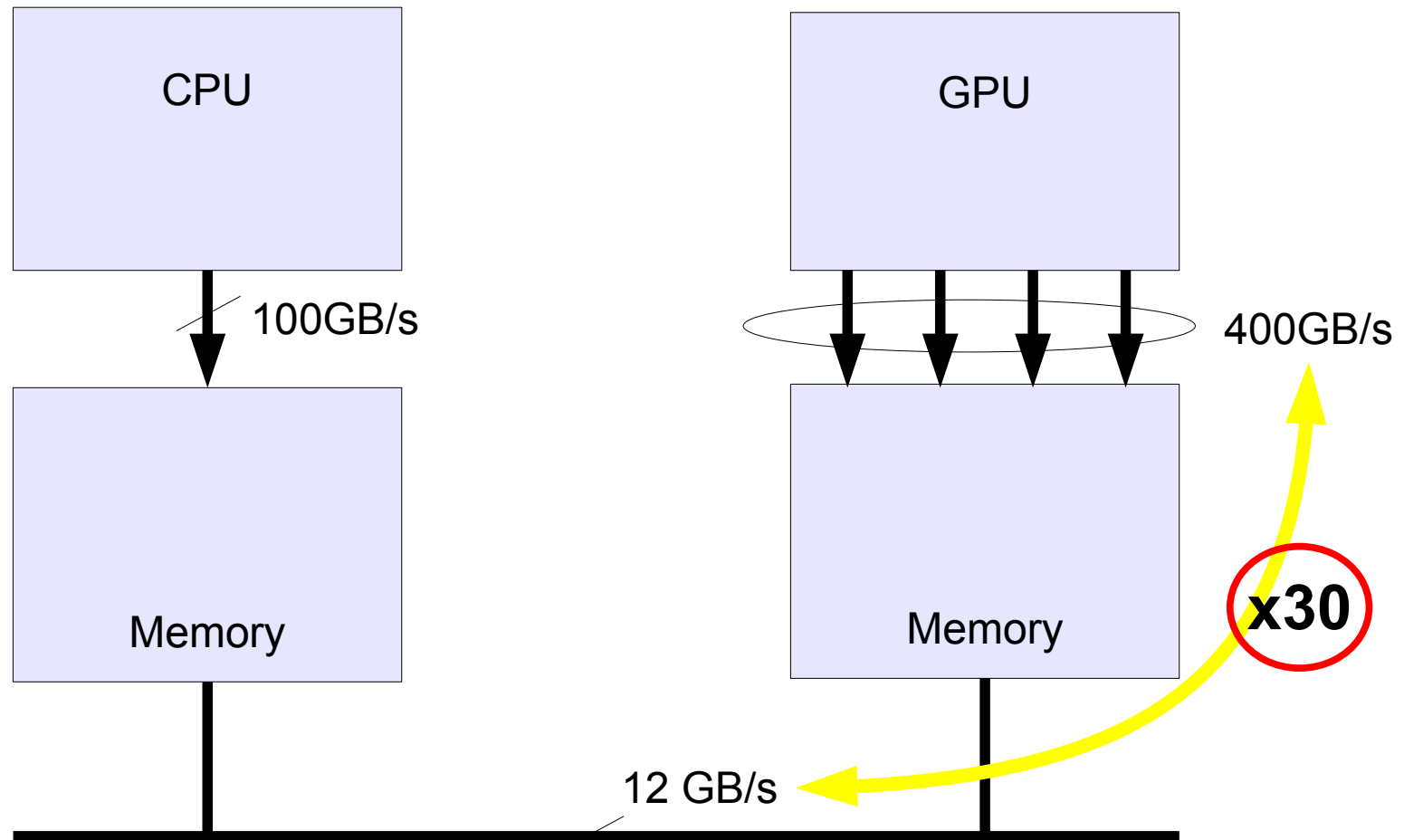
GPUs and data locality



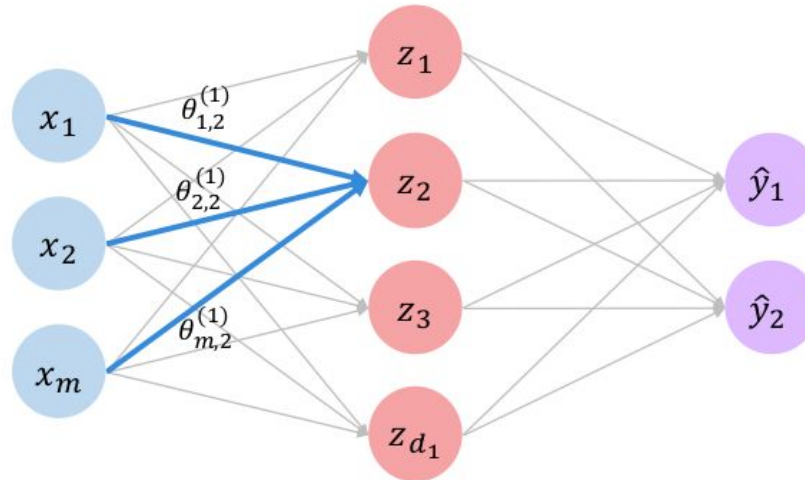
GPUs and data locality



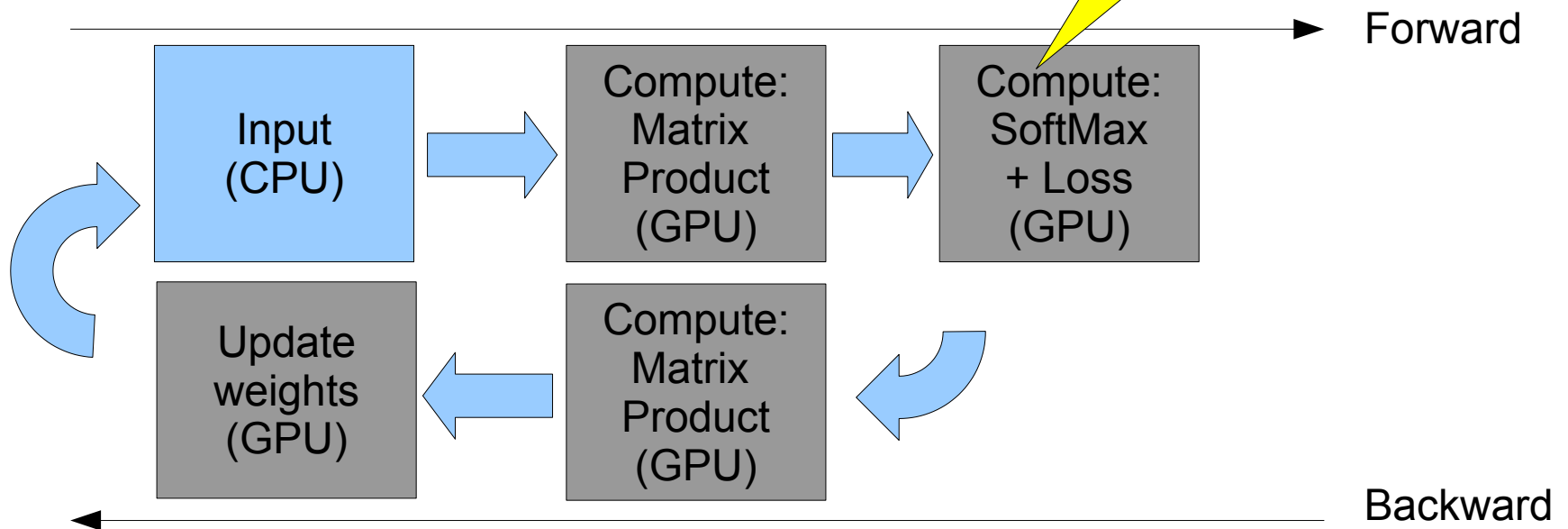
Problem: separate GPU memory



Data must be on a GPU

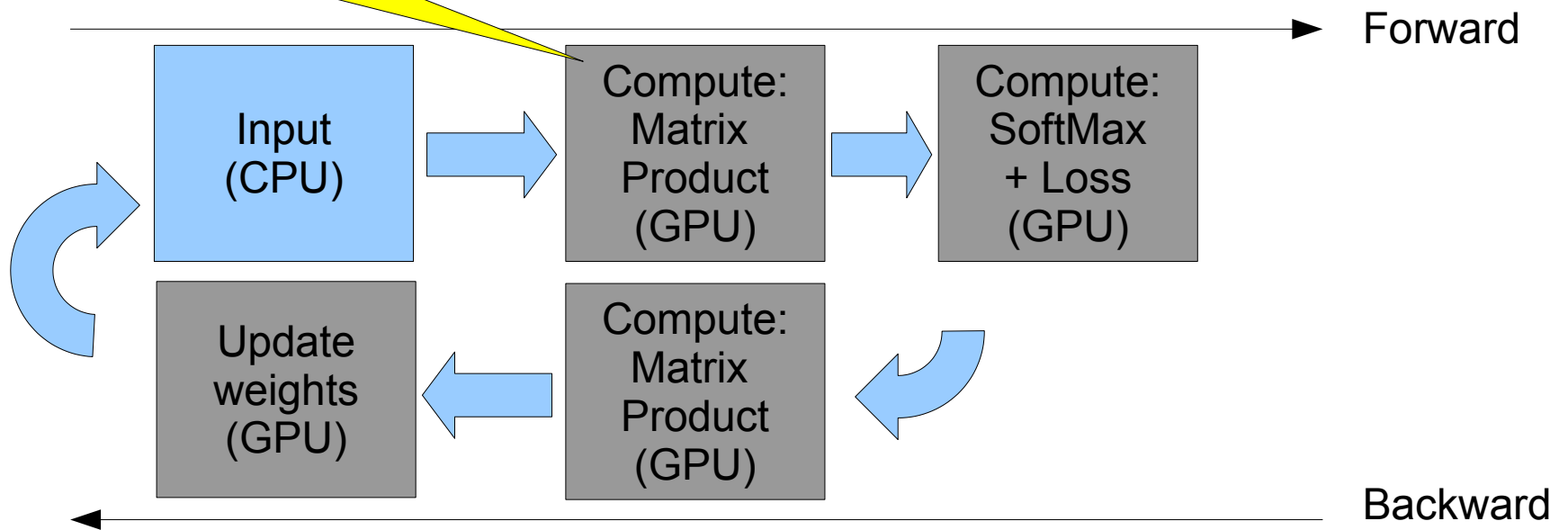
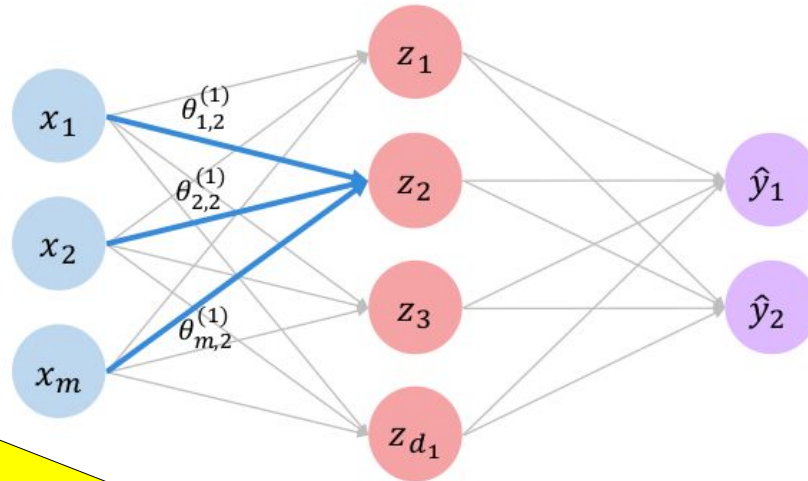


Cost(GPU)=
Cost(compute)+
Cost(move data)

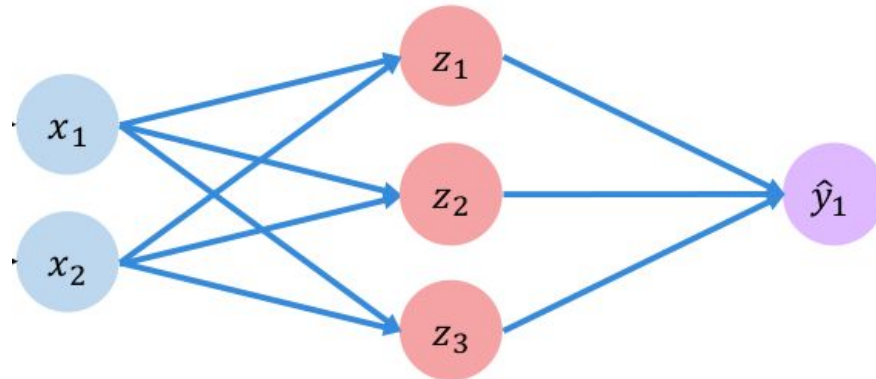


Data must stay on a GPU

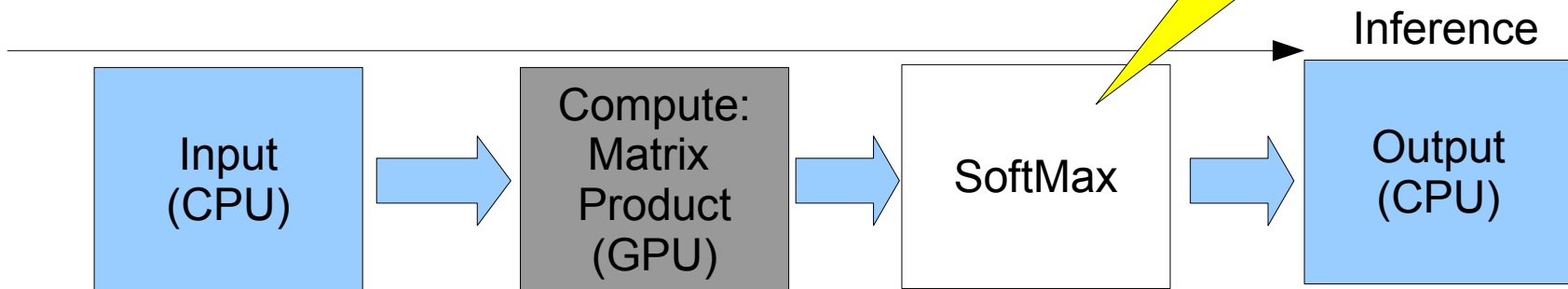
And what if data does not fit?



GPUs and invocation cost



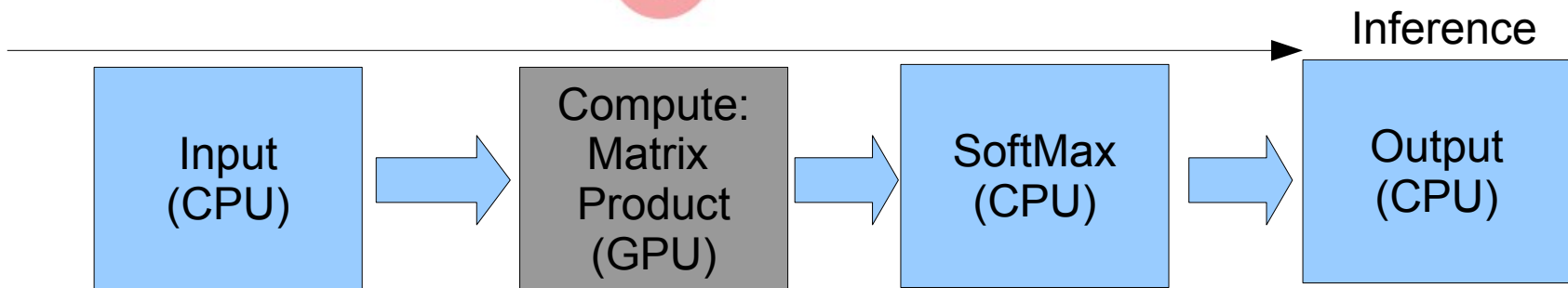
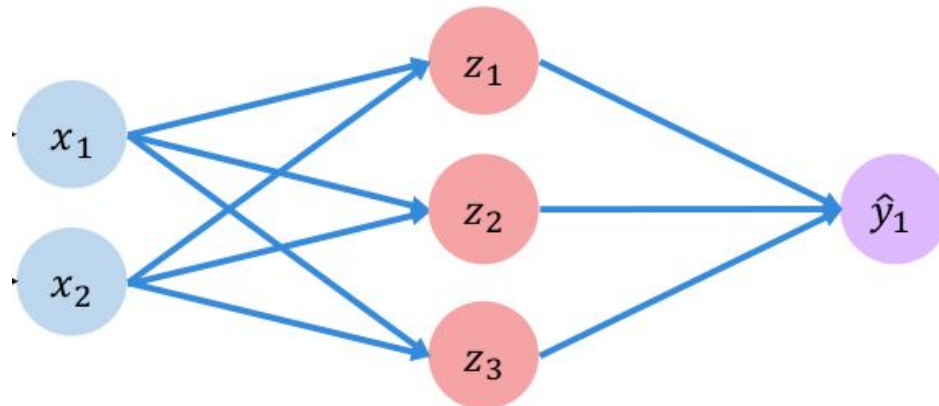
Where to compute?



Management overhead

- GPU invocation takes at least 30,000 single CPU core cycles
- Short GPU invocations do not pay off
 - Batch multiple invocations

Sometimes CPU is faster even when it's slower



if $\text{Cost}(\text{GPU}) > \text{Cost}(\text{CPU})$
=> compute on CPU

GPUs are used automatically, but..

- cuBLAS, cuDNN, cuFFD
 - provide powerful GPU-accelerated functions
- TensorFlow, MxNet, Caffe,...
 - run on GPUs
- So why learn how to use GPUs?

Other interesting projects in my lab

- Accelerator-centric systems
- Computer Networks with Accelerators
- Computational Storage with Accelerators
- Accelerator security
- Accelerators and OSes
- Integration of accelerators in Data Centers

Want to hear more?

- Accelerators and accelerated systems:
236/046-278
- Undergraduate/master projects @ Accelerated Computing Systems Lab (ACSL)

mark@ee.technion.ac.il

Zisapel 517, ECE

<https://acsl.group>